

TOFT seminar notes for presentation on 1 Mar understand ↴  
 simp. non-triv. eg. of TOFT what is going on in finite gp. version of Chem-Simons.  
2D TOFTs and Frobenius algebras fr. alg. perspective  
 & simp. non-triv. TOFT w/o Chem-Simons.  
 eg. of TOFT & analytic differential (Dijkgraaf-Witten  
 purpose perspective of path integral theory  
 } and understand an important  
 come together in a very nice way! first example of TOFT.

Today's aim: 'Clarify' all 2D TOFTs & understand an important  
 Recall last week: definition of TOFT + classification of 1-D TOFT.  
 Let's review definition of TOFT as certain things will become simp.

~~Def~~ of  $(n-1)+1$ -dim TOFT: (today,  $n=2$ )

Category Bord<sub>n</sub>: objects: closed, oriented  $(n-1)$ -folds / (orientat.-preserving)  
 bordisms: (oriented) bordisms  $M$  fr.  $E$  to  $F$  / (orientat.-pres.)  
 $\xrightarrow{E \rightsquigarrow F} \xrightarrow{E \rightsquigarrow F}$  diffeomorphisms  
 rel.  $E, F$

Convention:  $\{M\}$

Bord<sub>n</sub> has monoidal structure: obvious disjoint union.

+ symmetric (monoidal) structure:

recall in Vect<sub>F</sub>, symm. structure given by  $U \otimes V \xrightarrow{\sim} V \otimes U$

In Bord<sub>n</sub>,



: NOT (in general) identity even if  $U=V$ !

(recall equiv. of morphisms must hold  $U \otimes V$  and  $V \otimes U$  fixed)

TOFT is symm. monoidal functor fr. Bord<sub>n</sub> → Vect<sub>F</sub>.

such that  $\square \mapsto \otimes$   
 such braiding to braiding!

(orientat.-preserv)  
 diffeomorph to  $\square$

→ objects: pts<sup>1-D</sup>.  
 + 2 different objects -

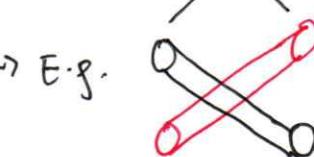
2-D

$\square$ . So ~~all~~ "all" objects are  
 $\emptyset, \square, \square \square, \dots$

Upshot: since we identify  $\square$  and  $\square$ ,  
 there's no need to keep track of orientat's today (or rather, all orientat's are  
 implicit).

Morphisms:

identified as same object.



Is this  $\simeq$  ? Ans: NO! As we have to hold source + target fixed: the color coding makes this clear (cannot send black to red).

However: TOFT being a symm. functor means it must always send



to  $(U \otimes V \rightarrow V \otimes U)$  Hence why we usually do not think much of it.

Recall Aim 1: classify 2D TFTs. So let  $\mathcal{F}$  denote any 2D TFT.  
 Since we focused on duality last week, today we will take an approach  
 that emphasizes that.

In particular recall last week: because we had a morphism that  
 looks like this:  , we saw:

(1)  $\mathcal{F}(0)$  is f.d., say =  $V$ .

(2)  $\mathcal{F}(0)$  is its dual  $\rightsquigarrow$  what we mean is  $\mathcal{F}(\overset{\circ}{0})$  gives a  
 non-degenerate pairing  $\langle \cdot, \cdot \rangle : V \times V \rightarrow \mathbb{F}$ .

(3) In fact, from the proof last week (recall we had things like

$$\text{2} = \underline{0} \underline{0}, \text{etc}$$

we can deduce  $\mathcal{F}(\overset{\circ}{0})$  must be the coevaluation map of the  
 pairing  $\langle \cdot, \cdot \rangle$ ; i.e. it sends  $1 \mapsto \sum V_i \otimes V_i^*$ .

In 1-D case, by considering , it is the end of the story.

Q: what can we do in 2-D that we can't in 1-D?

a) Recall last week we also saw this interesting feature of TFT:

If we have a cobordism  , then  $\mathcal{F}(0)$  essentially picks out  
 $\text{an element of } V$ .  
 (Can't do this in 1-D!)

b) Prototypical example when we say cobordism: 

$\mathcal{F}(\overset{\circ}{\square})$  is a map  $V \otimes V \rightarrow V$  (Can't do this in 1-D!).

This is nothing but a multiplication in  $V$ , making  $V$  an  $\mathbb{F}$ -algebra!  
 check: - unit? Is precisely the ele given by  $\mathcal{O}$ , because of

$$\text{3} = \underline{0} \underline{0} \quad *$$

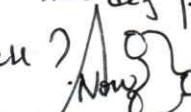
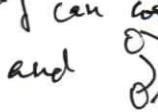
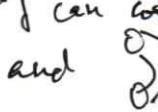
$$\text{---} " \sim \otimes 1 \rightarrow \sim "$$

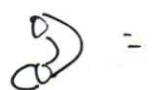
- associativity? Yes, as

$$\text{4} = \text{3} \text{3} !$$

so now: given 1D TQFT  $\mathcal{F}$ , get  $V$  a f.d.  $\mathbb{F}$ -alg (unital, associative)

Non-deg pair can ~~have~~ with a non-deg. pair  $\langle \cdot, \cdot \rangle$ .

Is that all?  and  look similar. Let's try to apply the maneuvers of unital & associativity to .

- Unital:  =  $0$ , a linear functional  $V \rightarrow \mathbb{F}$   
defined by  $v \mapsto \langle v, 1 \rangle$ .

Nasty work, ~~but this will~~ for now  
(but we will return to it later!)

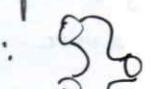
- Associativity:

$$\text{Diagram} = \text{Diagram} \quad \text{i.e. } \langle xy, z \rangle : \langle x, yz \rangle.$$

Such  $\langle \cdot, \cdot \rangle$  is also called associative!

Upshot: given 2D TQFT  $\mathcal{F}$ , get  $V$  a f.d.  $\mathbb{F}$ -alg (unital, assoc.)  
with non-deg associative pair  $\langle \cdot, \cdot \rangle$

Def. V a Frobenius alg.!

One of the most important properties in algebra: commutativity → comm. alg.  
Is  $V$  comm.? Yes!! Check:  = 

$$\begin{matrix} "x" \\ "y" \end{matrix} \mapsto xy \quad \begin{matrix} "x" \\ "y" \end{matrix} \mapsto yx$$

Recall need hold source and target fixed. But we can just twist!

Have shown:

Prop: Given 2D TQFT  $\mathcal{F}$ , get a comm. Frob. alg.  $V$ .

In fact:

Then. The converse is true: a comm. frob. alg  $V$  is sufficient to fully

define a 2D TQFT  $\mathcal{F}$ . Forget about having TQFT. Just start in  $V$ .

How to prove this? First qn: what is needed to define a 2D TQFT?

Recall for 1D TQFT: look at generators & relations.

For 2D TQFT; we can do the same b/c we have classification of surfaces.

Prop. Bord $_2$  is generated by object  $\emptyset$  (under disjoint union) } ele.  
and morphisms  (and 

Rule: What does it mean to 'generate'? Every morphism can be written as a composite/disjoint union where in each piece, just  $\neq$  one of the above 'elementary' morphisms.

$$\begin{array}{c} \emptyset \quad b \\ \emptyset \quad b \\ \emptyset \quad b \end{array}$$

Pf' Morse theory (indices of critical points corr. to the above) .

Rule: Why can't we do the same in 3D and higher? Actually, clue is in the above use of Morse theory: a critical pt. of index 1 is just a saddle: how do we know it is locally a pair of pants (we can already see there are two diff. directions, so this is not spherical)?

This requires a detailed argument which is only feasible in the 2D case (cf. Hirsch, Diff. Top) essentially the pt of claspers of surfaces

This also hints at the need for extended TFT: If we allow just a saddle by itself to count as a morphism, then such obstacles would be reduced. But that requires us to also allow 0-dim mfd's in a 2D TFT: so much is precisely what extended TFT does?

So need to define TFT on these elementary basisns - We all know what

$\mathcal{O}$  and  $\mathcal{G}$  have to be: the unit & multiplication. So suff. consider  $\mathcal{O}_S$  and  $\mathcal{O}$ . (Also we earlier said  $\mathcal{G}$  is also known.)

Now we have all seen  $\mathcal{O}$  but deferred it to later! ( $\mathcal{G}$ ,  $\mathcal{O}$ ) by another

It must be the linear flat given by  $v \mapsto \langle v, 1 \rangle$

Degression: There is 1-1 corr. between assoc. pairs and linear functionals.

Given assoc. pair, we have seen how to produce a linear functional.

Conversely, given linear functional  $\mathcal{O}$ , from " $\mathcal{G} = \mathcal{O} \mathcal{O}$ " we see that a pair should be given by  $\langle x, y \rangle = \mathcal{O}(xy)$ .

The non-deg condition corr. to:  $\ker \mathcal{O}$  has no (non-trivial) left ideals.

So can also define Tch. alg. by existence of such linear flat: more useful algebraically, but the non-deg condn has no good interpretation in TFT setting.]

More interesting one is  $\mathcal{O}_S$ . Recall that for  $\mathcal{G}$  we compared it with  $\mathcal{G}$ .

So now we compare it with  $\mathcal{G}$ ; recall at the start we said this must be the ~~assoc~~ coevaluation corr. to the non-deg. pair  $\mathcal{O}$ , i.e.  $1 \mapsto \sum v_i \otimes v_i^*$ .

How to define  $\mathcal{O}_S$  using  $\mathcal{G}$  and what we ~~are~~ currently have?

Start w : now need one input  
 outputs  $\vee$  so should combine two into one:  
 (but want two) =

and now have three  
 & what we are saying:  
 If we have such a  $V \rightarrow V \otimes V$   
 which can be used to define  
 a TFT, it has to be equal  
 to this composition. Same for  
 $V \rightarrow \sum v_i \otimes v_i^* \rightarrow \sum (vv_i) \otimes v_i^*$ .

In other words  $\Rightarrow$  it must correspond to  $V \rightarrow V \otimes V$  sending

Obvious qu: does this coincide with ? If it doesn't, then there is no hope of defining a TFT starting with  $V$ . This is sth that has to be checked independently of TFT considerations, i.e. purely algebraically:

Is  $\sum (vv_i) \otimes v_i^* = \sum v_i \otimes (v_i^* v)$ ?

Standard technique: have non-deg form  $\langle \cdot, \cdot \rangle$ , its basis & dual basis.

Multiplicit:  $vv_i$  is hard to handle. So write

$$\sum_{i,j} \langle vv_i, v_j^* \rangle v_j \otimes v_i^* \quad \sum_{i,j} v_i \otimes v_j^* \cancel{\langle v_j, v_i^* v \rangle} \quad \sum_{i,j} v_i \otimes v_j^* \langle vv_j, v_i^* v \rangle$$

Now swap  $i, j$ .

So get a 'canonical'  $V \rightarrow V \otimes V$  starting fr.  $V$  a comm. Frob. alg.  
 This is a commutant on  $V$ !

Similar to multiplicit: check:

① Comut: a  $V \rightarrow V$ , obvious candidate is !

want sth like " =

i.e.  $\sum \langle vv_i, 1 \rangle v_i^* = v$ ? Yes!

② Coassociativity: " =

$$v \mapsto \sum (vv_i^*) \otimes v_i \rightarrow \sum_{i,j} v_j \otimes (v_j^* (vv_i^*) \otimes v_i)$$

$$v \mapsto \sum v_j \otimes (v_j^* v) \rightarrow \sum_{i,j} v_j \otimes ((v_j^* v) v_i^*) \otimes v_i$$

③ Cocommutativity: " =

$$v \mapsto \sum (vv_i) \otimes v_i^* \rightarrow \sum v_i^* \otimes (vv_i)$$

$$v \mapsto \sum v_i^* \otimes (v_i v)$$

so essentially follows from coassociativity.

so essentially follows from commutativity.

Rule: Now  $V$  has unit, mult, comunit, counit which ~~satisfy~~ correspond to the four elementary braidings to be used to define the 2D TFT.

Rule that this does not make  $V$  a bialgebra, as comunit, <sup>the here</sup> sends  $1 \mapsto \sum v_i \otimes v_i^*$  (where in bialg. comunit sends  $1 \mapsto 1 \otimes 1$ .)

Although some  $V$  bialgs (e.g. Hopf alg.) are also then Frob. alg..  
impt

This means we have two different comunits on the alg.

Now we have specified ~~what the~~ what map the TFT should send each ele. cobordism to.

Now we need to check the most imp't thing: is this well-defined??.

i.e. given any ~~two~~ cobordism and two diff. decompositions into ele.

cobordisms, do the con. linear maps agree?

Sketch pf.: show that they coincide with the lin. map attached to a certain cobordism, both in normal form by relations which we know hold between the lin. maps assigned to the ele. cobordism;

# idea: Given two decoups, put them both in normal form by relations which we know hold between the lin. maps assigned to the ele. cobordism;

These two normal forms must coincide by genus consider., etc.

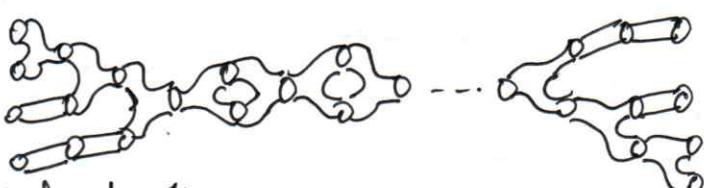
Steps: ① Do above for conn. cobordisms w/o braid ~~crosses~~  $\rightarrow$  this is the meat of the proof!

② For conn. w/ braid ~~crosses~~, inductively remove them (not immediate but just technicality)

③ For disconnected cobordism: just add some ~~two~~ braids at the

end to view them as disjoint unions ( $\sqcup$ )

Normal form:



(if source =  $\emptyset$ , then start w/  $\textcircled{1}$ ;  
similar for target).

① How to braid to normal form? Move all  $\textcircled{1}$  left (all  $\textcircled{2}$  right.)

Along the way of moving left, what can  $\textcircled{1}$  meet? (resp.  $\textcircled{2}$  right)

- Identity: pass through  $\textcircled{1} \textcircled{2} = \textcircled{2} \textcircled{1}$   
( $\text{id} \circ \text{mult} = \text{mult} \circ \text{id}$ )

-  $\textcircled{1}$ : disappear:  $\textcircled{1} \textcircled{2} = \textcircled{2} \textcircled{0}$

(in  $V$ , the unit is indeed unit for mult.)

-  $\textcircled{2}$ : two cases.

(resp.  $\textcircled{2} \textcircled{1} \dots$ )

(counit)

a)  : then just view it as middle part of normal form.

b)  This represents a composition

$$v \otimes w \mapsto \underbrace{\sum v_i \otimes (v_i^* v) \otimes w}_{\text{which is hence equal to}} \underbrace{\sum v_i \otimes (v_i^* v w)}$$

which is hence equal to 

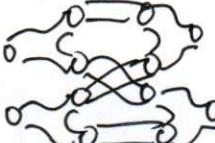
Similarly

 = 

This is the most imp relation, called the Frobenius relation, and  $\Rightarrow$ : we have exhausted all cases, also the last of the generating relations we need.

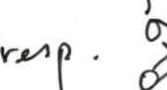
link: This relation is about "how multiplex<sup>2</sup> & comultiplex commute".

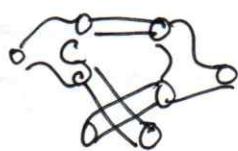
In a bialgebra, comut is an alg. hom., so mult & comut

commute by  = 

Step

② To handle  : can assume inductively surrounding ~~areas~~ areas are in normal form. So only a few cases :

 (resp.  ) : wocomm. resp. comm.

 :  $v \otimes w \mapsto \sum (vv_i) \otimes v_i^* \otimes w$   
 $\mapsto \sum (vv_i) \otimes w \otimes v_i^*$   
 $\mapsto \sum (vwv_i) \otimes v_i^*$

so it is equal to  again, allowing us to remove this braidy.

This completes the proof (sketch) of the main theorem. 

Remarks on pt of the main thm:

Pt is ~~long~~! interesting but brief: because it can't generalise!

~~Facts on the key takeaways: none in~~

However, it can help us answer a more interesting qn:  
as always, when we want to study other, we should always start in the simplest case & see what it tells us.

What if we want to generalise this to higher dimensions?

what works:

1-D: specify on  $\circ^+$ .

2-D: specify on  $\circ$ .

what doesn't:

3-D ~~sphere~~ (say): specify on all surfaces of genus  $g \in \{0, 1, 2\}$ ??

Fix: extended TFT: extend all the way down to O-fields (pt) and specify on pt.

"Cobordism hypothesis": suffice to specify on pt!"

Morse theory to split into elementary cob.

Duality: 1-D: tells you  $\circ^-$   
2-D: non-deg. part -

Morse theory to split into elementary cobordisms.

~~Fix~~: Cobordism hypothesis: use duality data in essential way.  
Notion of "fully dualizable"

E.g. in 2-D case: need (much more) extra work to say that saddle pt is just pair of pants.

Fix: extended TFT to split into more local cobordisms (e.g. just saddle itself). This also gives us more flexibility when using it to compute invariants (next).

In particular: above treatment of TFTs requires us to understand all surfaces completely. ~~Still~~, so using it to then compute invariants attached to surface is somewhat 'reverse':

What we will do next: illustrate how TFT is usually used to compute invariants: find a TFT defined 'by nature' i.e. 'globally', then see what it does locally and use that to compute.

### Example

(1 other similar thg e.g. Hept alg)

Most standard example of Frobenius alg. from algebraist's perspective:  
group alg.  $C[G]$  for a finite gp  $G$ .

(char of alg - closed)

of course, this is not comm. Take its center  $V$ : those elements  
with same coefficient on each conjugacy class.  $\cong$  class funct's on  
 $G$  (in product as convolution).

Rep theory of finite gps: characters of reps are class fcts, and there  
is a bilinear form

$$\langle \phi, \psi \rangle = \frac{1}{|G|} \sum_{\phi_1, \dots, \phi_k} \phi(g) \psi(g^{-1})$$

wrt which the class of the inreps of  $G$  form an orthonormal basis  
(b/c this form essentially computes the dim of the Hom space)

Can check this pairing is assoc., so  $V$  is ~~rep~~ a Frob alg. w/ this  
pairing.

The unit  $1 = \sum_{\phi} \phi(1) \phi$ .

and conv. lin. fctal  $\phi \mapsto \langle \phi, 1 \rangle = \frac{\phi(1)}{|G|}$  ( $\cong \dim V_\phi$ ) where  $V_\phi$   
(count)

( $= \dim V_\phi$  if  $\phi$  is char of inrep.)

Since have pnb  
let the class of the inreps be  $\phi_1, \dots, \phi_k$ . ( $k = \# \text{conj. classes. of } G$ )

Last thing we don't know abt  $V$  is its comultiplication.

Suppose it sends  $\phi \mapsto \Delta(\phi) = \sum_{j,k} c_{jk} \phi_j \otimes \phi_k$ .

Then  $\Delta$  has to satisfy:  $\phi_i = \sum_{j,k} c_{jk} \phi_j \otimes \phi_k = \sum_{j,k} c_{jk} \phi_j \otimes \phi_i$

Comult is simple:  $v \mapsto \{v\phi_i\} \otimes \phi_i$ .

and counit?

" $\delta$ "

So Need to understand the multiplication in  $V$  wrt.  $\phi_1, \dots, \phi_k$ .

In fact one can show  $\psi_i = \frac{\phi_i(1)}{|G|} \phi_i$  ~~satisfy~~  $\psi_i^2 = \psi_i$ ,  $\psi_i \psi_j = 0$ .

(why? The  $\psi_i$  act as scalars on each inrep  $\phi$ , taking trace & using the  
orthogonality relation we see they act as id on its con. inrep and  
0 on the others, so they are the primitive central idempotents in the  
block decompr. of  $C[G]$ .)

So  $\phi_i^2 = \left(\frac{\phi_i(1)}{|G|}\right)^{-1} \phi_i$ ,  $\phi_i \phi_j = 0$ .

By ~~note~~  $V$  defines a TFT  $\mathcal{T}$ , and

~~Now want to use  $\mathcal{T}$  to compute invariants of surfaces.~~

Closed  
Surface  $\Sigma_g$  of genus  $g$ : compute  $\chi(\Sigma_g)$  locally in normal form:



~~$\int \phi_i d\sigma$~~

each of these sends

$$v \mapsto \sum (v\phi_i) \phi_i \mapsto v(I\phi_i^2)$$

~~$\int \phi_i d\sigma$~~   $\mapsto v(I\phi_i^2)$

so

$$1 \longmapsto \sum \phi_i^{2g} \quad (\text{recall } \phi_i \phi_j = 0)$$

$$\sum \left( \frac{\phi_i(1)}{161} \right)^{2g} \phi_i \mapsto \sum \left( \frac{\phi_i(1)}{161} \right)^{2-2g} = \sum \left( \frac{\dim V_i}{161} \right)$$

$\Sigma_g$   
 $\chi(\Sigma_g)$

Observe: only invariant of  $\Sigma_g$  we get is just its Euler characteristic  $\chi(\Sigma_g)$ .

This whole process hinges on fact that we know clstrat<sup>1</sup> of surfaces;

It cannot give us any interesting invariants beyond those which we already know <sup>a priori</sup>. clstrat<sup>1</sup> of surfaces (i.e. genus,  $\chi$ , etc.)

~~To have hope of~~ Also: above const. is artificial & essentially rep-theory.

~~Does it arise more naturally?~~ Soln: We need to define the TQFT by ~~as~~ a process independently of the known clstrat<sup>1</sup> of surfaces. Then we can check that the local pieces agree with the local def<sup>1</sup> above and the computat<sup>1</sup> of the invariant will be by chopping into pieces as above.

Recall from last week the "motivat" for TQFT".

To each object in Bord<sub>n</sub> ( $(n-1)$ -fold) is assigned a Hilbert space which is a space of funct's on some kind of space of fields.

So here suppose for  $E$  a 1-fold we ~~also~~ have "space of fields" = principal G-bundles over  $E$   $L^2(\mathrm{Princ}(E))$  and the cov. Hilbert space is funct's on iso. classes of pin. G-bundles over  $E$ .

To a cobordism should con. a linear map between the two f<sup>1</sup> spaces which should be given by a sort of path integral.

so need measure on "space of fields"

For  $P \in \text{Prin}_G(E)$ , define  $\mu(P) = \frac{1}{|\text{Aut}(P)|}$ . Why? This gives the so-called groupoid cardinality! Given cobordism 

want lin. fct  $L^2(\overline{\text{Prin}_G(E)}) \rightarrow L^2(\overline{\text{Prin}_G(F)})$

Have diagram

$$\begin{array}{ccc} & \phi^* & \\ \text{pullback} \swarrow & & \downarrow \text{restrict} \\ \text{Prin}_G(M) & & \text{Prin}_G(F) \\ \text{restrict} \swarrow & & \downarrow \text{restrict} \\ \text{Prin}_G(E) & & \end{array}$$

cf. Terence Tao's blog post, and in fact what we will do here has been studied systematically by Baez-Hoffnung-Walker as "degroupoidification"

How to pushforward?

Integration along fibers.

New function  $\gamma$  should send  $y \in \overline{\text{Prin}_G(F)}$

$$\begin{array}{c} \gamma \mapsto \int_{\text{Prin}_G(M)} \gamma(y) \mu(y) = \int_{\text{Prin}_G(M)} \phi^*(z) \mu(z) \\ \text{at } \gamma(z) \text{ s.t. } z \text{ restricts to } y \\ \text{"path integral" to } y \\ \text{which is just} \\ \text{a finite sum!} \end{array}$$

$$\begin{aligned} \gamma(y) &= |\text{Aut}(y)| \sum_{z \in \text{Prin}_G(M)} \phi^*(z) \frac{1}{|\text{Aut}(z)|} \\ &\quad z \text{ restricts to } y \\ &= |\text{Aut}(y)| \sum_z \phi(z|_E) \frac{1}{|\text{Aut}(z)|} \end{aligned}$$

In fin.-gp. case, no analytic difficulties, everything is rigorous.

Here the action functional is trivial.

Can check this is indeed functorial (respects compo<sup>s</sup>) (can refer to Bas paper above) so this defines a TFT independent of choice of surfaces!

Concretely, what does this mean, in our example, esp.?

Need to understand the main player "space of fields"  $\text{Prin}_G(E)$  better, with  $E=S^1$  as our <sup>discrete</sup> example.

First of all: since  $G$  ~~is finite~~,  $\text{Prin}_G(E)$ -bundle is just a (regular) covering space over  $E$  & not necessarily connected.

We know such covering spaces are classified by quotients of  $\pi_1(E)$  (its group of deck transformations).

e.g.  $E=S^1$  ( $G=\mathbb{Z}/3\mathbb{Z}$ )



In this way one can establish:

$$\underline{\text{Prin}_G(E)} \cong \underline{\text{Hom}(\pi_1(E), G)}.$$

$G$  connected

$$\text{Prop: } \forall \underline{\text{Prin}_G(E)} \cong \underline{\text{Hom}(\pi_1(E), G)}/\text{G-conj}.$$

Rules:

Deck(CP)

- ~~Deck(CP)~~ is a subgp of  $G$ , but need not be whole  $G$ .

- ~~A local conn~~ Principal  $G$ -bundle has extra structure

( $G$ -action), so consider maps  $\underline{\text{Hom}(\pi_1(E), G)}$

In the language of fibres, (i.e. embeddings of Deck(CP) into  $G$ ) here, essentially implicitly explained, the transport is via holonomy / monodromy have a choice of base pt.

since  $G$  discrete there is a natural flat connect<sup>on</sup> on any bundle; giving notion of "parallel transport": locally we have  $\underline{\underline{\text{---}}}$

which allows us to consider holonomy: given an ele. of  $\pi_1(E)$ , going once around the con. loop we may end up at a diff. pt., giving a con. ele. of  $G$ . (essentially manifested by the fact it is con. sp.)

When are two bundles equivalent? given a loop in  $E$ , may lift it to  $P$  - morphisms of  $G$ -bundles will conjugate the holonomy. (Example for  $G = \mathbb{Z}/3\mathbb{Z}$ )

Principal  $G$ -bundles have more structure than cover spaces, so this approach is rather low-tech; still a few details to check + make it all precise.

A more high-brow proof: recall  $\underline{\text{Prin}_G(E)} \cong [E, BG]$ .

In fact for  $G$  discrete, w/ long exact homotopy seq on  $\underline{\underline{E}}$  and fact that characterizing prop. of univ. bundle is  $\underline{\underline{E}}$  has all trivial homotopy gps (weakly contractible), we see that  $BG$  is a  $K(G, 1)$ .

Then  $\underline{\underline{E}}$ , for a base pt  $e \in E$ ,  $[E, BG]_{\text{base pt}} \cong \underline{\text{Hom}(\pi_1(E), G)}$

(PT in chap 1 to Hatcher's alg. top textbook & essentially obstruct-theoretic)

so  $[E, BG] \cong \underline{\text{Hom}(\pi_1(E), G)}/\text{G-conj}$ .

Back to our example.

We see  $\tau(O) = f^*$ 's on conjugacy classes of  $G = \mathbb{V}$

Multiplication =  $f(O) \cdot g(O) = h(O)$  given two  $\neq$  class  $f^*$ 's f.g.

the corresponding funct<sup>n</sup>  $h$  is obtained by

To want to look at 'all'  $\text{Prin}_G(E)$   
i.e. with a fixed basepoint  
then  $\text{Prin}_G(E) \cong \underline{\text{Hom}(\pi_1(E), G)}$

$$h_{\text{cpl}} = |\text{Aut}(y)| \sum_{z \in \text{Princ}_G(\mathcal{Z})} (f \otimes g)(z|_0) \frac{1}{|\text{Aut}(z)|}$$

t restricts  
to  $y$

~~$y \in G/G\text{-conj.}$~~

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$$\text{Princ}_G(\mathcal{O}) = \sum_{y_1, y_2} |C_G(y_1)| \sum_{\substack{(y_1, y_2) \in G \times G \\ \text{not } G\text{-conj.}}} f(y_1) g(y_2) \frac{1}{|C_G(y_1, y_2)|}$$

~~$y_1, y_2 \in G/G\text{-conj.}$~~   
 $y_1, y_2 \text{ conj. to } y$

$$\text{-stab.} = |C_G(y)| \sum_{\substack{(y_1, y_2) \in G \times G \\ y_1, y_2 \text{ conj. to } y}} f(y_1) g(y_2) \frac{1}{|G|}$$

$$\text{orb.-stab.} = \frac{1}{\#\text{conj. class}} \sum_y \sum_{\substack{(y_1, y_2) \in G \times G \\ y_1, y_2 \text{ conj. to } y}} f(y_1) g(y_2)$$

$$= \sum_{y_1, y_2 \in y} f(y_1) g(y_2) \quad (\because f, g \text{ are class fns})$$

which is indeed convolution!

Similarly one checks unit  $\mathcal{O}$  and  $\mathcal{O}$  pairing are as we defined them w.r.t. the Frobenius alg.  $V$  (they are ~~as~~ simpler & easier than the multiplication).

So this does indeed give the same TFT  $\mathcal{Z}$ .

However, what is the partition fn this time?

$$h(\text{genus surface}) = \sum_g \frac{1}{|\text{Aut}(z)|} \quad (\text{groupoid cardinality})$$

$z \in \text{Princ}_G(\Sigma_g)$

$$= \sum_{z \in \text{Hom}(\pi_1(\Sigma_g), G)} \frac{1}{|\text{Aut}(z)|}$$

$\text{not } G\text{-conj.}$

$$\text{-stab.} = \sum_{z \in \text{Hom}(\pi_1(\Sigma_g), G)} \frac{1}{|G|} = \frac{|\text{Hom}(\pi_1(\Sigma_g), G)|}{|G|}$$

which counts principal  $G$ -bundles on  $\Sigma_g$ .

$\mathcal{O}$  has  
 $\mathbb{Z} \cong \mathbb{Z} * \mathbb{Z}$   
 given princi. bundle,  
 holonomy around top leg  
 $= g_1$  and around  
 bottom leg  $= g_2$ ,  
 then around  
 one leg  $= gh$ .

Combining w prev., we see that we may compute this invariant by

$$\frac{|\text{Hom}(\pi_1(\Sigma_g), G)|}{|G|} = \sum \left( \frac{\dim V_i}{|G|} \right)^{\chi(\Sigma_g)} \quad \boxed{!}$$

For ~~more~~ complicated surfaces i.e. high  $g$ , LHS is hard to compute  
while RHS is very easy!

This is perhaps the simplest important example of computing an invariant (of closed surfaces) via TaFT.

Next week: go towards 3-dimensional manifolds.